

Waddle

Maintaining Canonical Form After Edge Deletion

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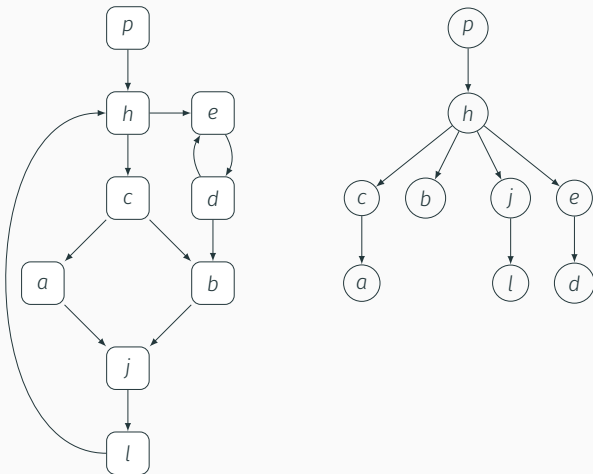
What Does Waddle Maintain? (1)

Dominator Tree
encodes which blocks occur on all paths to another block

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What Does Waddle Maintain? (2)

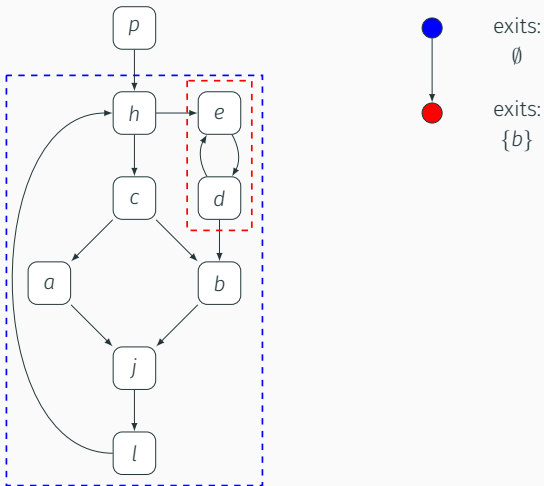
Loop Nesting Forest

encodes loop body sets · loop exit sets · loop nesting structure

What Does Waddle Maintain? (2)

Loop Nesting Forest

encodes loop body sets · loop exit sets · loop nesting structure



What Does Waddle Maintain? (3)

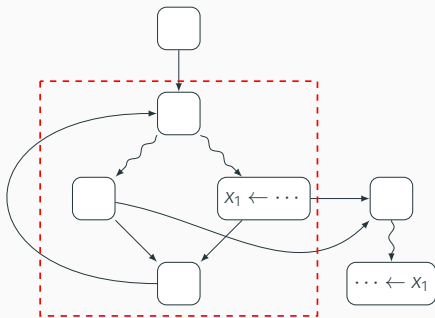
SSA + LCSSA Form

all names defined once · uses of name occur within defining loop

What Does Waddle Maintain? (3)

SSA + LCSSA Form

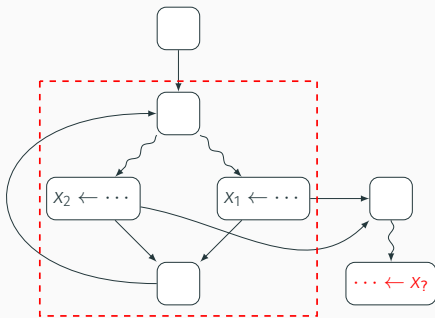
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What Does Waddle Maintain? (3)

SSA + LCSSA Form

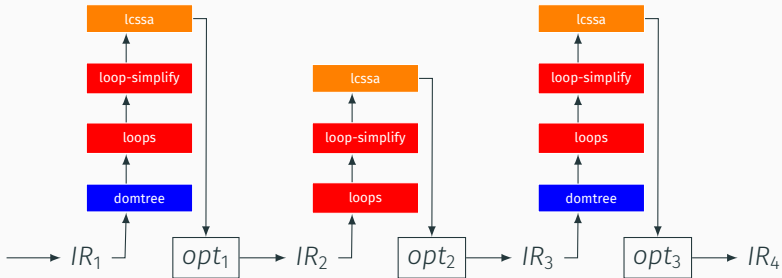
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'Canonical' Properties
LLVM's Loop Simplify Form

Optimization Pipeline Strategies

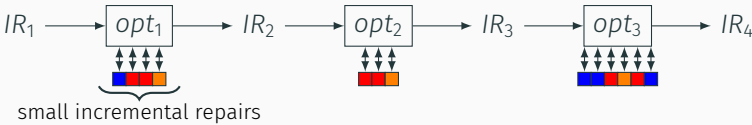
Repair On-Demand (LLVM's Approach)



LLVM 6.0.0 -O2 Passes

| | | | | | |
|------------------|----------------------|----------------------|----------------------|----------------------|----------------------|
| ipsccp | propagation | simplifycfg | lazy-value-info | basiccg | simplifycfg |
| globalopt | simplifycfg | domtree | jump-threading | globals-aa | domtree |
| domtree | domtree | aa | lazy-value-info | float2int | loops |
| mem2reg | aa | loops | correlated- | domtree | scalar-evolution |
| deadargelim | loops | lazy-block-freq | propagation | loops | aa |
| domtree | lazy-block-freq | instcombine | domtree | loop-simplify | demanded-bits |
| aa | instcombine | loop-simplify | aa | lcssa | lazy-block-freq |
| loops | lbcalls-shrinkwrap | lcssa | memdep | aa | slp-vectorizer |
| lazy-block-freq | loops | scalar-evolution | dse | scalar-evolution | instcombine |
| instcombine | branch-prob | indvars | loops | loop-rotate | loop-simplify |
| simplifycfg | block-freq | loop-idiom | loop-simplify | loop-accesses | lcssa |
| basiccg | lazy-block-freq | loop-deletion | lcssa | lazy-block-freq | scalar-evolution |
| globals-aa | pgo-memop-opt | loop-unroll | aa | loop-distribute | loop-unroll |
| prune-eh | domtree | mldst-motion | scalar-evolution | branch-prob | lazy-block-freq |
| inline | aa | aa | licm | block-freq | instcombine |
| functionattrs | loops | memdep | postdomtree | scalar-evolution | loop-simplify |
| domtree | lazy-block-freq | lazy-block-freq | adce | aa | lcssa |
| sroa | tailcallelim | gvn | simplifycfg | loop-accesses | scalar-evolution |
| aa | simplifycfg | aa | domtree | demanded-bits | licm |
| memoryssa | reassociate | memdep | aa | lazy-block-freq | globaldce |
| early-cse-memssa | domtree | memcpyopt | loops | loop-vectorize | constmerge |
| speculative- | loops | scop | lazy-block-freq | loop-simplify | domtree |
| execution | loop-simplify | domtree | instcombine | scalar-evolution | loops |
| domtree | lcssa | demanded-bits | barrier | aa | branch-prob |
| aa | aa | bdce | elim-avail-extern | loop-accesses | block-freq |
| lazy-value-info | scalar-evolution | aa | basiccg | loop-load-elim | loop-simplify |
| jump-threading | loop-rotate | loops | rpo-functionattrs | aa | lcssa |
| lazy-value-info | licm | lazy-block-freq | globalopt | lazy-block-freq | aa |
| correlated- | loop-unswitch | instcombine | globaldce | instcombine | scalar-evolution |

Always Canonical (Waddle's Approach)



```
// If we have a pass and a DominatorTree we should re-simplify impacted loops
// to ensure subsequent analyses can rely on this form. We want to simplify
// at least one layer outside of the loop that was unrolled so that any
// changes to the parent loop exposed by the unrolling are considered.
if (DT) {
    if (OuterL) {
        // OuterL includes all loops for which we can break loop-simplify, so
        // it's sufficient to simplify only it (it'll recursively simplify inner
        // loops too).

        // TODO: That potentially might be compile-time expensive. We should try
        // to fix the loop-simplified form incrementally.
        simplifyLoop(OuterL, DT, LI, SE, AC, PreserveLCSSA);
    } else {
        // Simplify loops for which we might've broken loop-simplify form.
        for (Loop *SubLoop : LoopsToSimplify)
            simplifyLoop(SubLoop, DT, LI, SE, AC, PreserveLCSSA);
    }
}
```

Canonical Form Loop Properties

Dedicated Preheader
enables easy + efficient instruction hoisting

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enables easy + efficient instruction hoisting

Dedicated Exit Blocks

enables easy + efficient effect sinking

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enables easy + efficient instruction hoisting

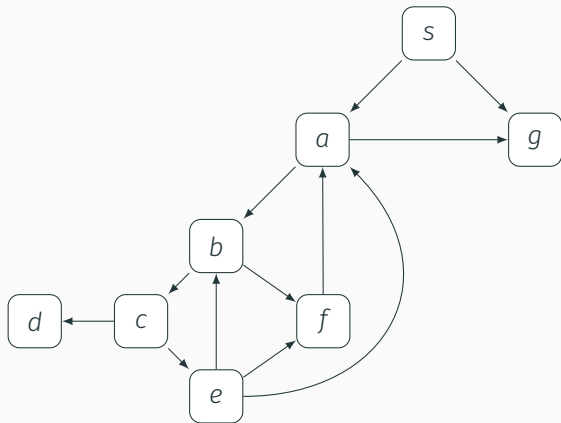
Dedicated Exit Blocks

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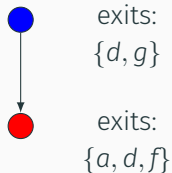
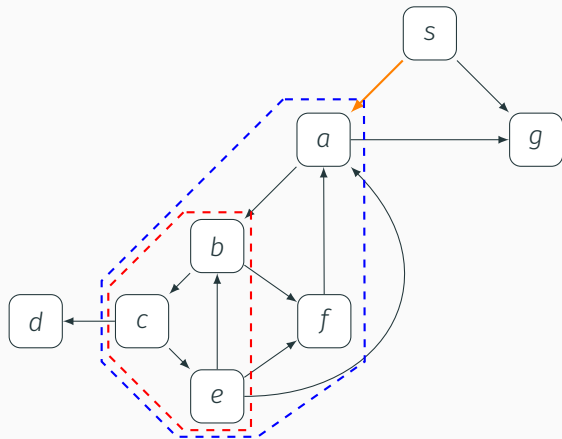
Unique Backedge + Latch

make destruction of loop unambiguous

Canonicalization

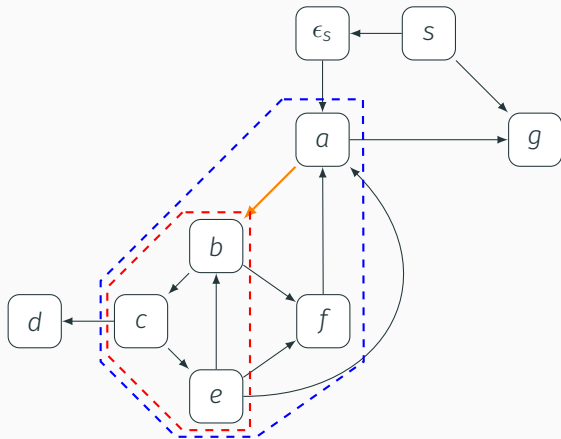


Canonicalization – Dedicate Preheaders



Construct loop nesting forest

Canonicalization – Dedicate Preheaders

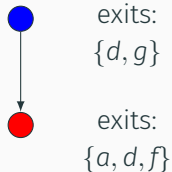
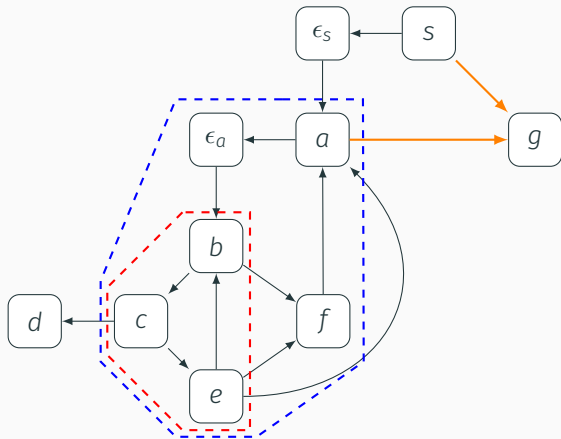


● exits:
 $\{d, g\}$

● exits:
 $\{a, d, f\}$

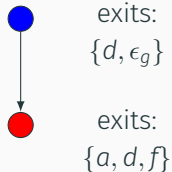
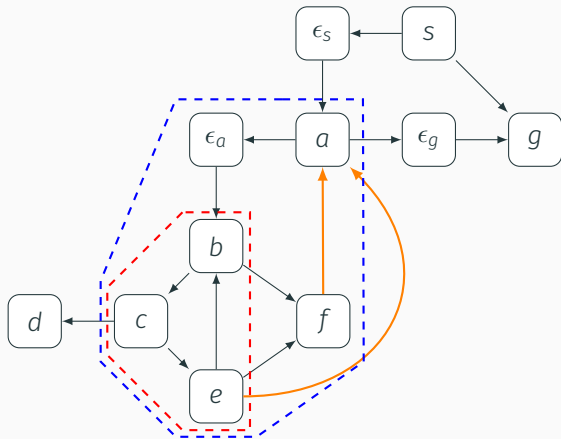
Dedicate preheader of outer (blue) loop

Canonicalization – Dedicate Exits



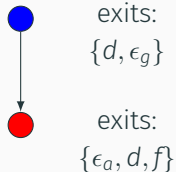
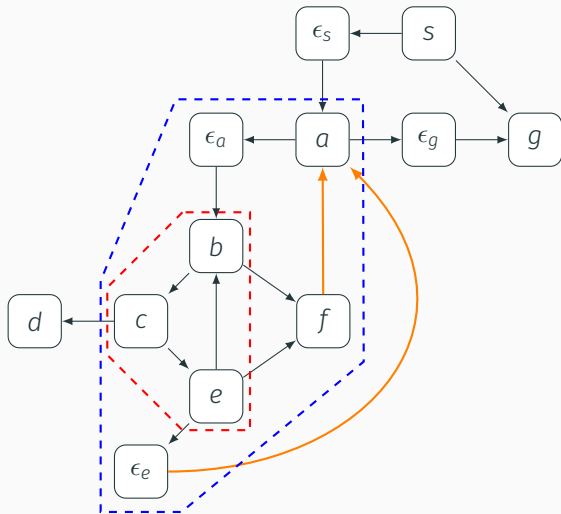
Dedicate preheader of inner (red) loop

Canonicalization – Dedicate Exits



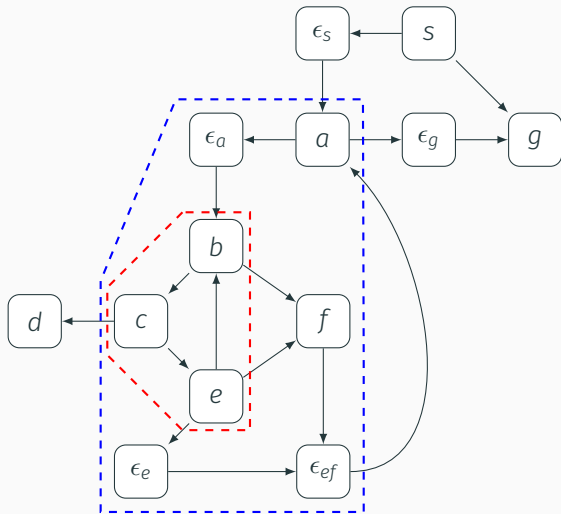
Dedicate exit (block g) of outer (blue) loop

Canonicalization – Ensure Unique Latches



Dedicate exit (block a) of inner (red) loop

Canonicalization – Ensure Unique Latches



exits:
 $\{d, \epsilon_g\}$

exits:
 $\{\epsilon_a, d, f\}$

Make latch for outer (blue) loop unique

Edge Deletion

(1) Remove edge from graph

Algorithm Outline

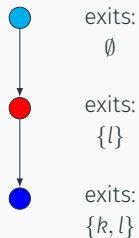
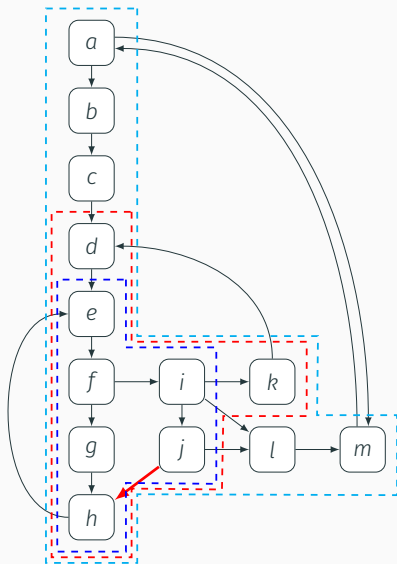
- (1) Remove edge from graph
- (2) Remove references to unreachable blocks and edges

Algorithm Outline

- (1) Remove edge from graph
- (2) Remove references to unreachable blocks and edges
- (3) Eject extraneous blocks from loop where edge was removed

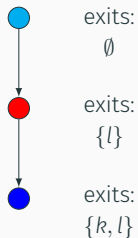
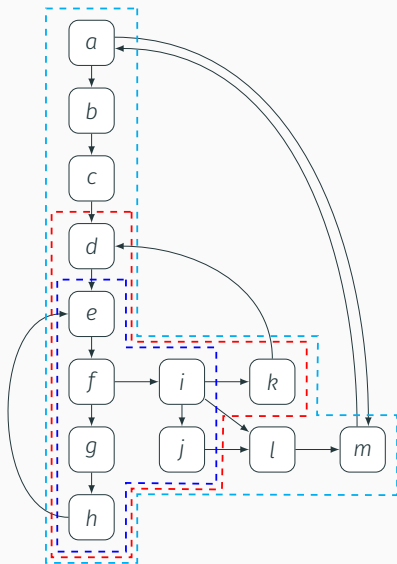
Example #1

Deleting Edge (j, h) – Incremental Repair



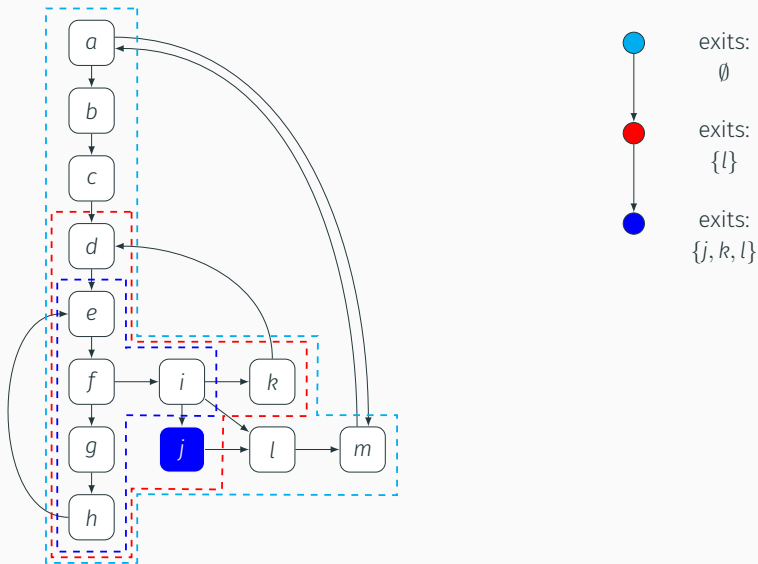
Initial graph

Deleting Edge (j, h) – Incremental Repair



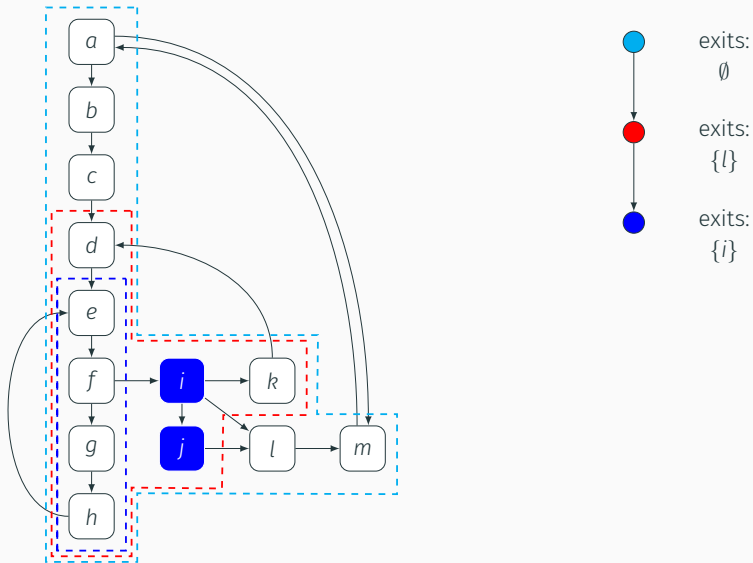
Edge deleted

Deleting Edge (j, h) – Incremental Repair



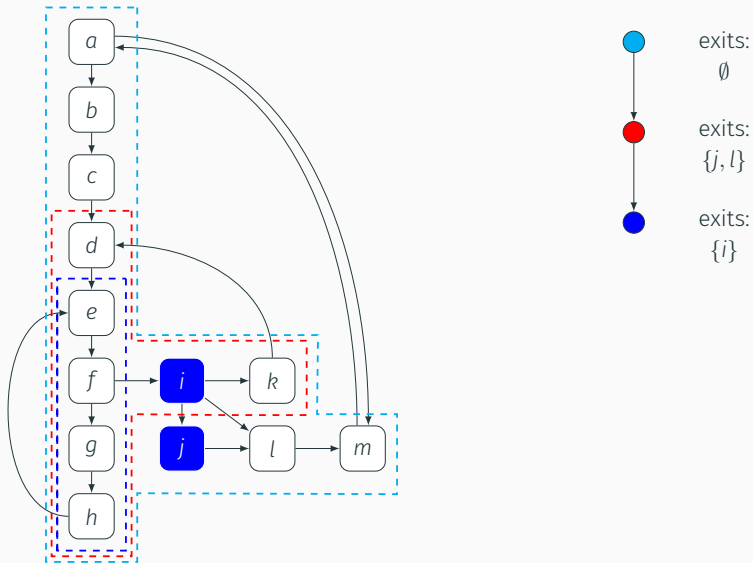
Eject block j from inner (blue) loop

Deleting Edge (j, h) – Incremental Repair



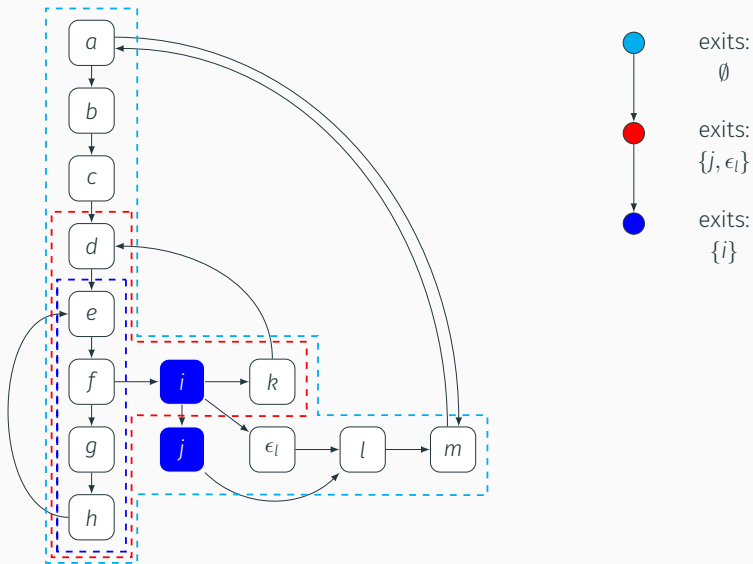
Eject block i from inner (blue) loop

Deleting Edge (j, h) – Incremental Repair



Eject block j from middle (red) loop

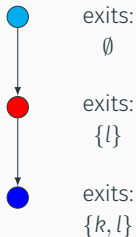
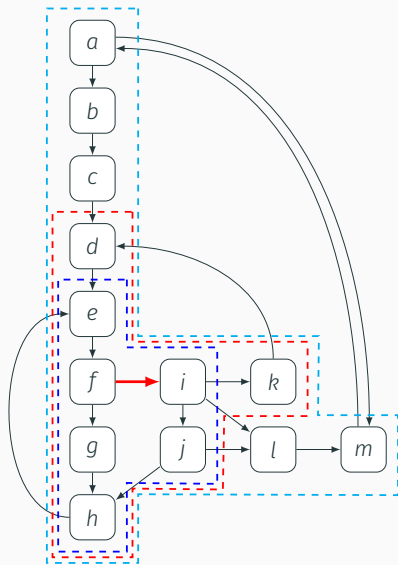
Deleting Edge (j, h) – Incremental Repair



Place block ϵ_l on edge (i, l) to dedicate exit

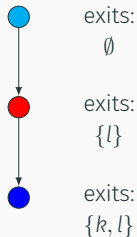
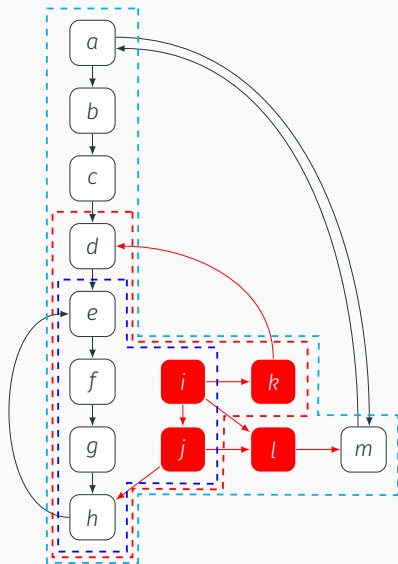
Example #2

Deleting Edge (f, i) – Incremental Repair



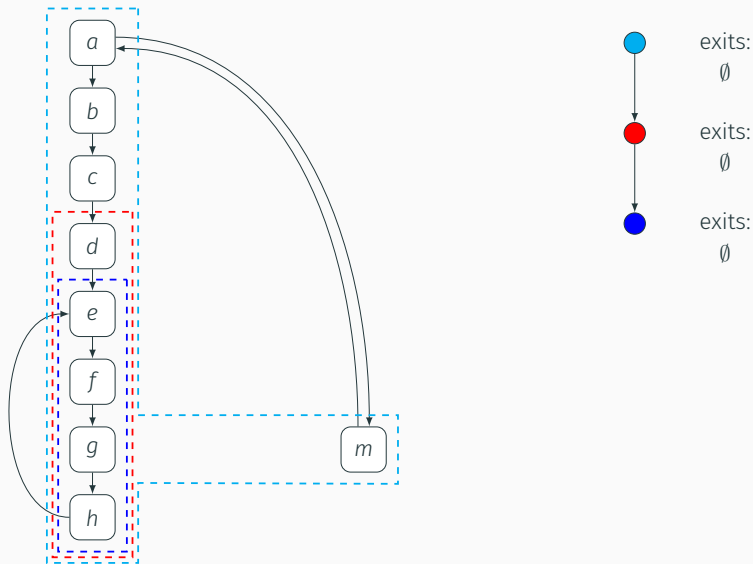
Initial graph

Deleting Edge (f, i) – Incremental Repair



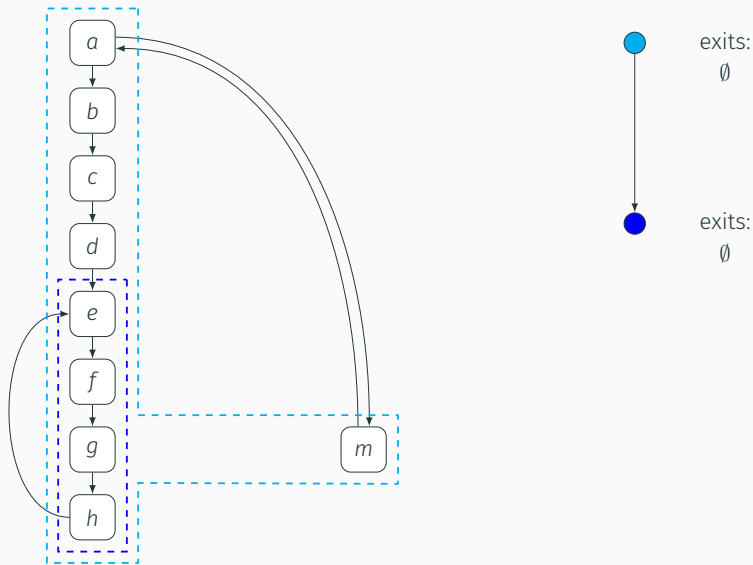
Edge deleted

Deleting Edge (f, i) – Incremental Repair



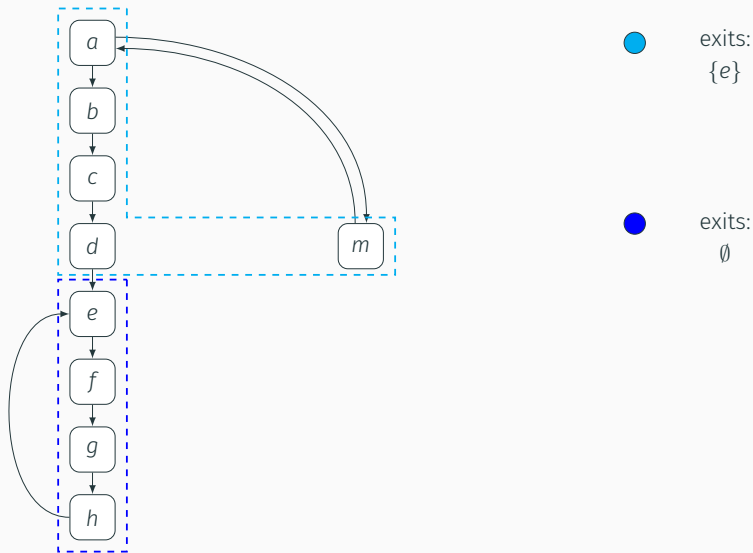
Remove unreachable blocks from graph, loop nesting forest

Deleting Edge (f, i) – Incremental Repair



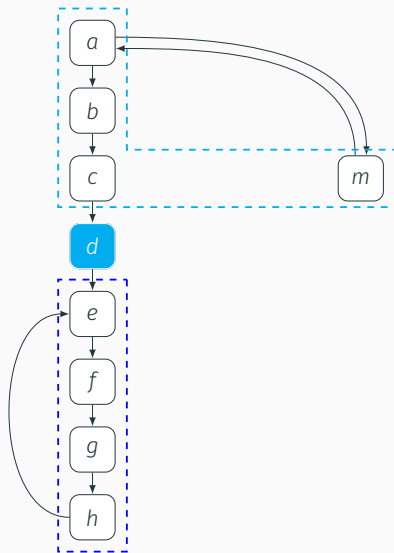
Remove destroyed middle (red) loop

Deleting Edge (f, i) – Incremental Repair



Eject block e (and its loop) from the outer (cyan) loop

Deleting Edge (f, i) – Incremental Repair

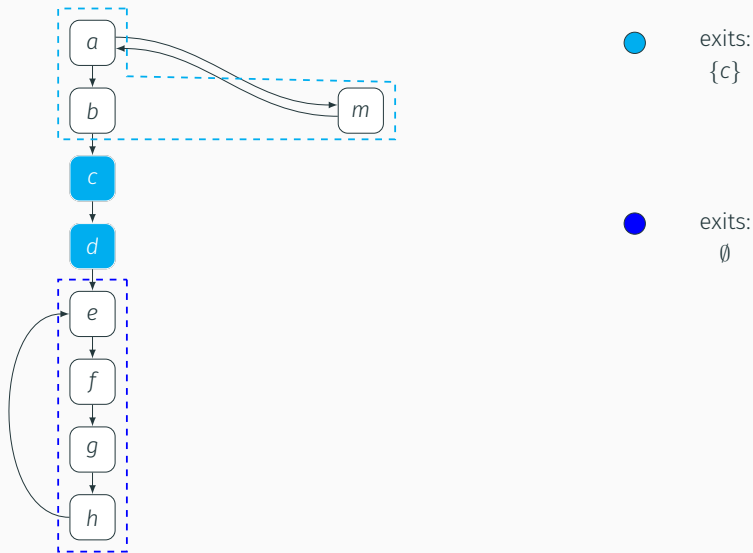


● exits:
 $\{d\}$

● exits:
 \emptyset

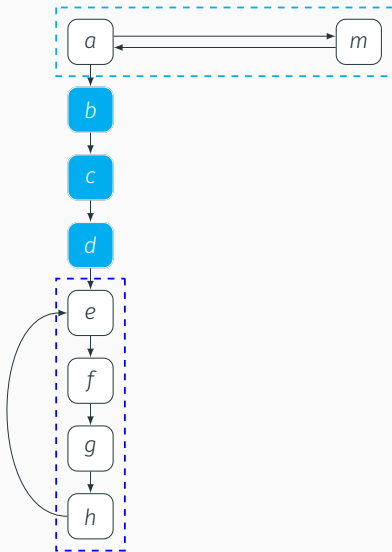
Eject block d from outer (cyan) loop

Deleting Edge (f, i) – Incremental Repair



Eject block c from outer (cyan) loop

Deleting Edge (f, i) – Incremental Repair



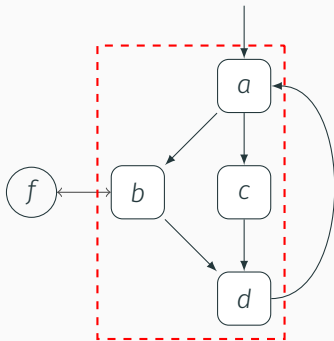
● exits:
{b}

● exits:
 \emptyset

Eject block b from outer (cyan) loop

Additional Applications

Function Inlining

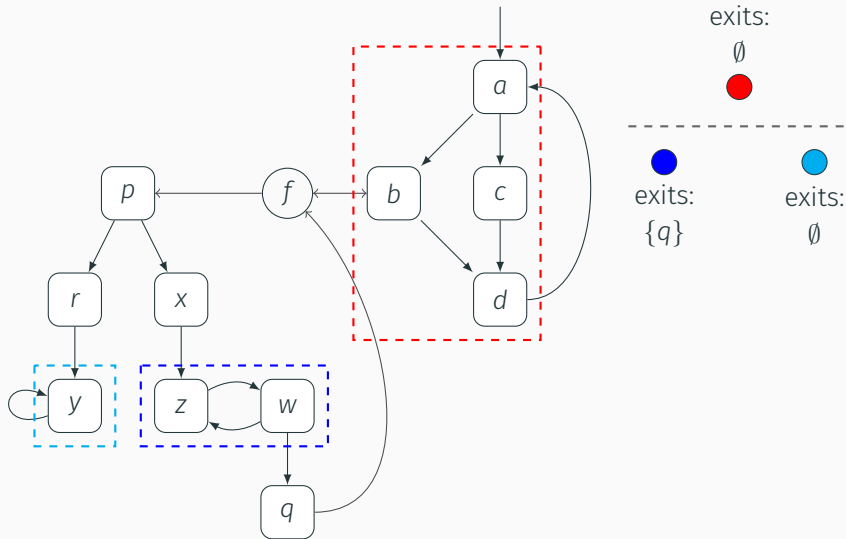


exits:

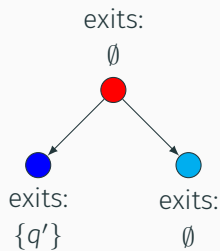
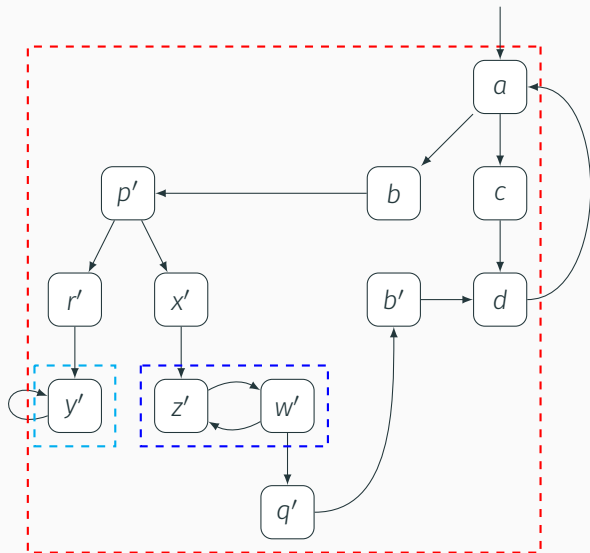
\emptyset



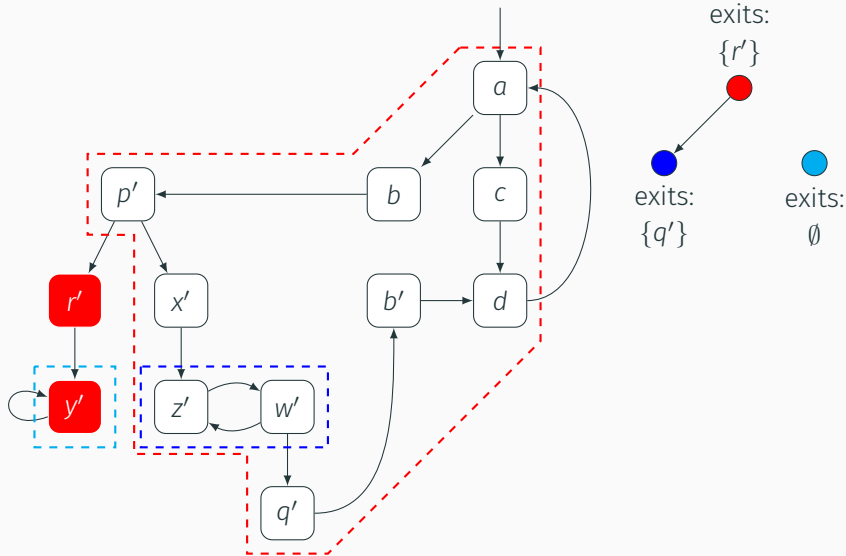
Function Inlining



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Function Inlining



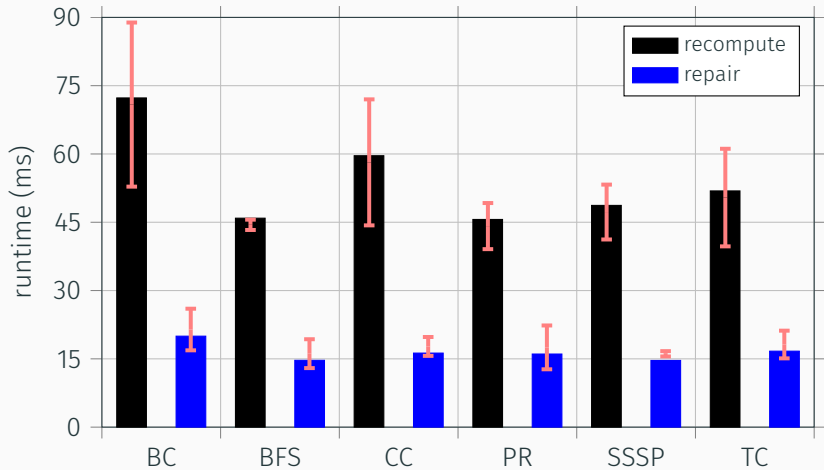
Evaluation

- (1) Construct Waddle IR from C++ source (through LLVM)
 - 6 compilation units
 - ~85 interesting functions per compilation unit
 - ~21 blocks, ~30 edges, ≤ 10 loops (\leq depth of 4) per function

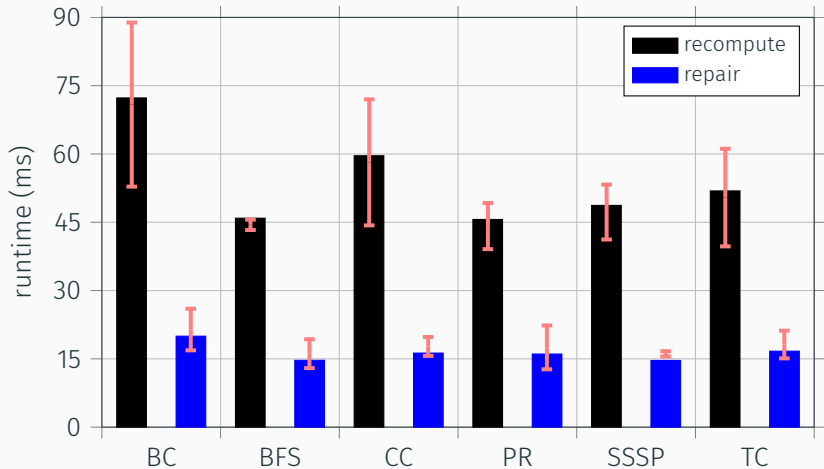
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 - ~85 interesting functions per compilation unit
 - ~21 blocks, ~30 edges, ≤ 10 loops (\leq depth of 4) per function
- (2) Construct a stable order of edges
- (3) For each edge that has siblings remaining:
 - Delete edge and reconstruct canonical form (baseline)
 - Delete edge using procedure described here

Results

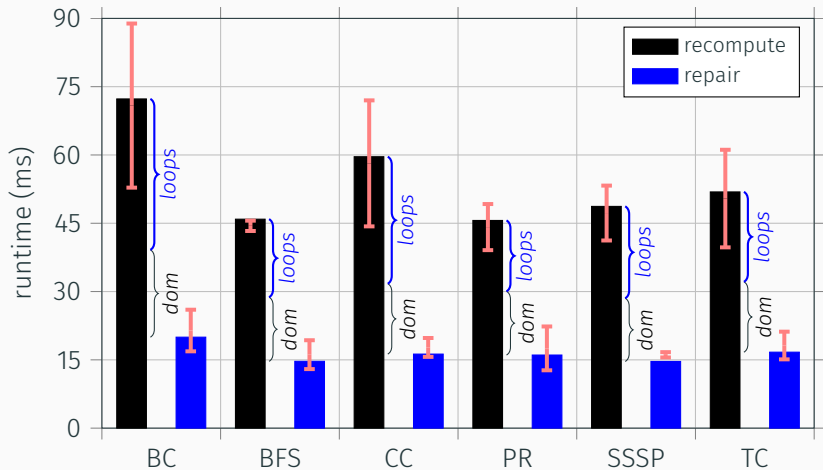


Results



64.80 to 72.7% decrease in runtime

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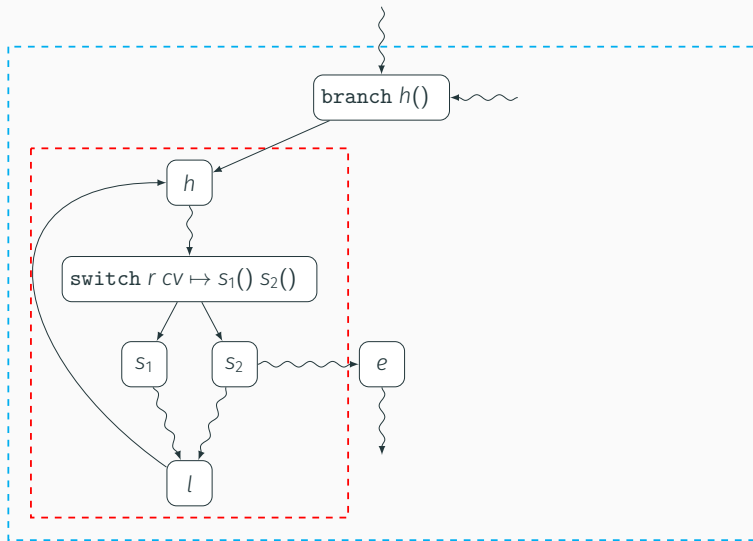


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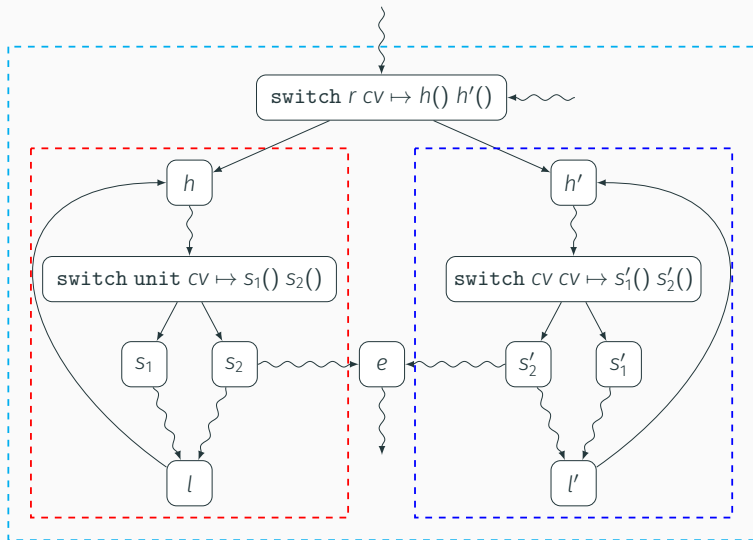
Questions?

Bonus Slides

Subgraph Duplication (Loop Unswitching)



Subgraph Duplication (Loop Unswitching)



Thank You!